Symbolic-Model-Guided Fuzzing of Cryptographic Protocols

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TL;DR. Critical and widely used cryptographic protocols have repeatedly been found to be flawed in their design and their implementation. A prominent class of such vulnerabilities are logical attacks, *i.e.,* attacks that solely exploit flawed protocol logic. Automated formal verification methods, based on the Dolev-Yao (DY) attacker, excel in finding such flaws, but operate on abstract specification models only. Fully automated verification of deployed implementations is today still out of reach. This leaves open whether the implementations actually being used are secure. On the opposite side of the spectrum, *fuzz testing*, developed since the 90s is now the gold standard for testing security software and is used at scale by the largest software companies. However, even if a protocol implementation is tested against memory-related vulnerabilities using state-of-the-art fuzzers, the whole class of *implementation-level logical attacks* remains out of scope. Unfortunately, this blind spot hides numerous attacks, notably recent logical attacks on widely used TLS implementations introduced by implementation bugs.

In prior work, we answer this with a novel approach combining domain-specific, security-related formal models (*e.g.,* DY attacker) and fuzz testing. We have designed and implemented a first Proof of Concept fuzzer written in Rust[1], which found four new vulnerabilities in the TLS library WolfSSL[2]. This recent work has opened up various exciting new research questions we would like to explore in this internship. In particular, by building on this prior work, an interesting project would to design and implement a DY fuzzer feedback metric, another one would be to reflect the varieties of DY security properties in the fuzzer engine.

A good command in Rust is recommended to start this internship.

Context. Today's information society crucially relies on secure information exchanges achieved by *cryptographic protocols*. Those distributed programs that leverage cryptographic primitives (*e.g.,* encryption, digital signature) to achieve various security goals are critical to many aspects of our modern society: finance, business, communication, etc. Any flaw in these protocols can have dramatic consequences, amplified by their ubiquity and our dependence on them. Yet, critical and widely used protocols have been repeatedly found to be **flawed in their design or their implementation**. A prominent class of such flaws are **logical attacks**, *i.e.,* attacks that solely exploit *flawed protocol logic* such as *Man-in-the-Middle* (MiM), *replay*, or *downgrade* attacks, etc.

As such flaws are subtle and hard to catch, **formal methods** have been proposed to analyze protocol *design specifications* since the 80s. Symbolic verification is a first-class, extremely successful such method [1]. It offers a **mathematical model capturing logical attacks**, *i.e.,* the

[1] Paper under submission, tool available at `https://github.com/tlspuffin/tlspuffin`

[2] CVE-2022-42905 (critical severity), CVE-2022-42905 and CVE-2022-42905 (medium severity), and CVE-2022-38153 (low severity).

symbolic model also called *Dolev-Yao model*, as well as rigorous and mechanized methods to reason about protocols. However, a fundamental, inherent issue is that symbolic verification **operates on abstract specification models only**. Security proofs thereon are of no practical use when the programs that end-users deploy or run are insecure. Unfortunately, history shows that frequent implementation bugs actually introduce vulnerabilities that were nonexistent in the specification, notably ***implementation-level logical attacks***. This is particularly well illustrated by the long history of such attacks in the ubiquitous and critical TLS and WiFi protocols ([2, 3, 8, 6, 7] to only name a few).

Programmers or auditors interested in precluding logical attacks from protocol implementations are left with ***testing*** since security-oriented *program verification* is extremely expertise-demanding and does not really scale beyond primitives or minimal protocols. As opposed to formal verification, testing is unsound by design (*i.e.,* bugs may be missed) but provides a certain level of confidence by excluding all the potential flaws covered by the body of tests. Therefore, a good *coverage* of the tests is paramount. Narrowing down to security, the gold standard is ***fuzz testing*** [4, 5] due to its ability to automatically generate test cases that maximize the coverage typically thanks to *feedback-driven evolutionary algorithms* utilizing *mutations*. Today, fuzzing is paramount in the **industry software development practices**, *e.g.,* Google, Cisco, Microsoft use it at scale. The state-of-the-art fuzzing techniques are adequate to find safety vulnerabilities (sometimes with potential security implications) but are unfortunately **unable to find logical attacks** since they operate at a too low level (*e.g.,* random bit-flips on network packets, code-based coverage). Prior works have proposed fuzzers operating in some ad-hoc *state machine model* [3, 2] that is also **too weak to capture the class of logical attacks**; *e.g.,* message contents cannot be tampered with by the adversary while most logical attacks rely on this.

Symbolic verification captures logical attacks at the design level only while fuzzing is industry-ready and operates on implementations but is limited to low-level, safety-oriented flaws, which are often the low-hanging fruits. Therefore, **effective and usable techniques to preclude logical attacks on implementations** are desperately lacking.

Objectives. This internship objective is to develop a **symbolic-model-guided fuzzing** framework. that will enable **checking implementations for the absence of logical attacks** with TLS as a case study. The central idea is to consider symbolic traces (from a symbolic model) as the input space of the Program Under Test (PUT) that will be fuzzed and then executed on the PUT through *concretization* (symbolic terms are evaluated into bitstrings). Fortunately, we already have a preliminary, proof-of-concept design and implementation in Rust for TLS 1.2 and 1.3 that will serve for this internship as a solid basis and test-bed for exploring new directions.

Intern's tasks. First of all, the intern will get familiar with formal verification in the symbolic model, fuzzing, as well as with the existing Rust code base.

Next, depending on the intern's affinity for and knowledge of the different involved aspects of this project, we will be able to adapt the project goals and choose one among several research directions, such as:

1. design domain-specific feedback metric to incentives the fuzzer to seek for new symbolic traces. The underlying fundamental question is: what is a good "symbolic feedback" that promotes semantically different symbolic traces?

2. define scoring metrics that can be effectively computed by symbolic verifiers and that can help the fuzzer promoting test cases that are close to attack traces.

3. design new fuzzing mutations and benchmark them with our test-bed,

4. design an efficient grammar-based fuzzing engine and evaluate it with our test-bed,

The precise direction this project will take shall be agreed upon with the intern at the beginning of the project. It can be more theory-oriented (1) or more practical-oriented (3-5). Should we find any vulnerability, we would follow standard and ethical responsible disclosure practices.

Expected ability of the student. We expect mathematical maturity, basic knowledge in logic, basic theoretical computer science. Knowledge in security and cryptography is not mandatory but is definitely a plus. For the implementation, a good command of Rust is recommended.

References

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