# Progresses in the analysis of Stochastic 2D cellular automata: <br> Asynchronous 2D Minority 

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# Asynchronous Systems 

Most of the real systems are asynchronous

- Networks, physical particles, biological cells...
- How does randomness introduced by asynchonicity affect the global behavior of these systems?


## Example A ring network

- where each node has two states: "has a token" or "does not have a token"
- running an algorithm that redistributes the tokens according to some rules/constraints:
"I get a token if none of my neighbors have one" or "I get a token if my right neighbor has one",...
- Example of question How long does it take to reach a "stable configuration"?


# Nature is an other example 



Patterns are governed by rule 30 as the shell grows

## 2D Cellular automata



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At each time step, each cell updates its states according to the state of its neighbors

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## 2D Cellular automata



Extensively used in physics, biology,... What happens in asynchronism regims?

## Cellular automata, here

- $0 / \mathrm{I}$ state ( $0=$ white \& I = black)
- Full asynchronism

A deamon chooses a random cell uniformly at random and updates it

- $\alpha$-asynchronism

Each cell is independently updated with probability $\mathbf{0}<\boldsymbol{\alpha}<\mathbf{I}$

Full synchronism: $\boldsymbol{\alpha}=\boldsymbol{I}$
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## Historic

Ergodicity of deterministic CA with random noise

- Toom - Gasc - Gray - Park - Louis (I974-)

Indecidability of independance to update history

- Gacs (2002)

Empiral studies of asynchronism

- Buvel, Ingerson - Bersini, Detour - Schönfish, de Roos (I995-)

Study of particular automata or classes of automata

- Fuks (2004 -)
- Fatès, Morvan, Regnault, S., Thierry (2005-)
- Chassaing, Gerin (2007)


# Historic: I D automata 6 types of relaxation times [FMST 2005, FRST 2006, CG 2007] 


(a) LOGARITHMIC (232)

(c') QUADRATIC (146)

(b) LINEAR (I30)

(d) EXPONENTIAL (2।0)

(c) QUADRATIC (I70)

(e) DIVERGING (I50)


## Fully asynchronous 2D Minority

- 0/I states
- $\mathrm{n} \times \mathrm{m}$ toric configurations
- A daemon selects uniformly at random a cell and the cell updates to the minority state among its 4 neighbors and itself.



## Energy function

Potential. $\mathbf{v}_{\mathbf{i j}}=\#\{$ neighbors in the same state as (i,j)\}

Cell (i,j) is active if $\mathbf{v}_{\mathbf{i j}} \geq 2$.
Energy of configuration c:
Energy $(\mathbf{c})=\sum_{(i, j)} \mathbf{v}_{\mathrm{ij}}$
Fact. The expected energy of a random configuration is
$2 \mathbf{N}$, where $\mathrm{N}=\mathrm{n} \mathrm{m}$

## Observed phase transition



## Observed phase transition



## Observed phase transition



## Energy is non-increasing when fully asynchronous

Theorem. The energy of a configuration is a nonincreasing function of time in fully asynchronous dynamic.

Energy decreases by at least 4 each time a cell with potential at least 3 is updated.
proof. $\Delta$ Energy $=8-4 \mathbf{v}_{\mathbf{i j}}$, when $(\mathrm{i}, \mathrm{j})$ is updated.

## Initial energy drop

Theorem. The energy of any configuration of size $N$ is at most $N+2 N / 3$ after $O\left(N^{2}\right)$ updates on expectation.
proof. Any such configuration contains a neighborhood in which a finite sequence of updates decreases the energy by at least I.

## Borders



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There is a red border between two neighboring cells in the same state.


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Fact. The borders are the boundaries of the covering homogeneous checkerboards regions

## Dual configurations

 (from now on, $n$ and $m$ are even)dual configuration $=$ configuration XOR checkerboard
$P+P$

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dual configuration $=$ configuration $\times$ OR checkerboard


## Stable configurations

Stable configurations are made of cells with potential $\leq 1$, i.e. in contact with at most one border:


Fact. Stable configurations are made of an even number of bands of width $\geq 2$ tiled by alternating checkerboards.


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## Energy of configurations



## Energy of configurations



# Typical asynchronous 2D minority 


$t=0$

$t=75 \mathrm{~N}$

$t=N$


$$
t=5 \mathrm{~N}
$$

$$
t=20 \mathrm{~N}
$$



$$
t=350 \mathrm{~N}
$$



$$
t=50 \mathrm{~N}
$$



$$
t=38 \mathrm{IN}
$$

# Convergence is almost sure 

Definition. The dynamics converges from an initial configuration $c^{0}$ if $\mathbf{T}=\min \left\{t: c^{t}\right.$ is stable $\}$ is almost surely finite.

Theorem. For all $c^{0}, E(T) \leq 2 N \cdot N 2 N$.
Proof. The following is a sequence of at most 2 N updates that stabilizes any configuration:
I. As long as there is an active black cell, flip it;
2. As long as there is an active white cell, flip it.

This sequence is followed with probability $1 / N^{2} \mathrm{~N}$.

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Definition. The dynamics converges from an initial configuration $c^{0}$ if $\mathbf{T}=\min \left\{t: c^{t}\right.$ is stable $\}$ is almost surely finite.

Theorem. For all $c^{0}, E(T) \leq 2 N \cdot N^{2 N}$.
Conjecture. If one of $n$ or $m$ is odd, the convergence time can indeed be exponential.

$2 n^{3}$

# Typical asynchronous 2D minority 


$t=0$

$t=N$


$$
t=5 N
$$


$t=20 \mathrm{~N}$


$$
t=38 \mathrm{IN}
$$

# Typical asynchronous 2D minority 


$t=0$

$t=N$

$t=5 \mathrm{~N}$

$t=20 \mathrm{~N}$


$$
t=50 \mathrm{~N}
$$

$$
t=75 \mathrm{~N}
$$


$t=150 \mathrm{~N}$

$t=300 \mathrm{~N}$

$t=350 \mathrm{~N}$
$t=38 \mathrm{IN}$

## Polynomial convergence

(conjecture. It is true as soon as $n$ and $m$ are even)
We study bounded configuration, which are surrounded by a checkerboard of width $\geq 2$ :


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# Typical evolution of a bounded configuration 


$t=N$


$$
t=20 \mathrm{~N}
$$


$t=5 \mathrm{~N}$

$t=25 \mathrm{~N}$

$t=10 \mathrm{~N}$

$t=30 \mathrm{~N}$

$t=15 \mathrm{~N}$


$$
t=35 N
$$

## Dual evolution:

## Automaton OT976


$t=N$
$t=5 \mathrm{~N}$

$t=15 \mathrm{~N}$

$t=20 \mathrm{~N}$

$t=30 \mathrm{~N}$

$t=35 \mathrm{~N}$

# Rules of the primal and dual dynamics 

Primal (Minority)

Dual
(Automaton OT976)


Isolated Peninsula



Corner


Inactive


Border Surrounded

## Coupling with the HV-convex hull


$t=N$

$t=20 \mathrm{~N}$

$t=5 \mathrm{~N}$

$t=25 \mathrm{~N}$

$t=10 \mathrm{~N}$

$t=30 \mathrm{~N}$

$t=15 \mathrm{~N}$

$t=35 \mathrm{~N}$

## Coupling with the HV-convex hull


$t=5 \mathrm{~N}$

$t=25 \mathrm{~N}$

$t=30 \mathrm{~N}$

$t=35 \mathrm{~N}$

## Rules of the hull dynamic


neighborhoods are considered with white/black symetries \& rotations (with exception of the bridge for the white/black symetry)

## Convergence

$$
\text { Let } \boldsymbol{f}=\text { \#black dual hull cells + Energy/4. }
$$

Lemma. For any island of the dual hull, $\mathrm{E}[\Delta f] \leq-3 / \mathrm{N}$.


## Convergence

$$
\text { Let } \boldsymbol{f}=\text { \#black dual cells + Energy/4. }
$$

Lemma. For any island, $\mathrm{E}[\Delta f] \leq-3 / \mathrm{N}$.
Proposition. For all bounded configuration,

$$
\begin{aligned}
\mathrm{E}[\Delta f] & \leq(2 \text { \#island contacts }-3 \text { \#islands }) / \mathrm{N} \\
& \leq \text {-\#islands/N. }
\end{aligned}
$$

Theorem. Every bounded configuration converges to the checkerboard configuration in finite time a.s.. The expected convergence time is:
$\mathrm{O}(\mathrm{N} \cdot$ Initial area).

## Conclusion

- An exponential upper bound on the convergence time.
- Polynomial convergence time for bounded configurations
(can be extended to configuration with a checkerboard band of width $\geq 2$ )
- An useful energy function that defines proper statistical physics
- Similar results for the Moore-Neighborhood


## Conjectures

- Phase transition at $\alpha_{c} \approx 0.83$.
- $\alpha_{c}=\frac{\sqrt[3]{46+\sqrt{6969}}}{6}-\frac{16}{3 \cdot \sqrt[3]{46+\sqrt{6969}}}+\frac{2}{3}$
- Some ideas to prove polynomial convergence time to checkerboard for $\alpha<$ I/3 but...
- Strongly depends on the underlying network (a lot of differences on trees...)
- Generalization to the class of threshold automata


## Thank you

