Dynamic of cyclic automata over \mathbb{Z}^2

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Cyclic Automata Network

- Graph G = (V, E)
- Set of states: $Q = \{0, 1, ..., q 1\}$
- Each vertex v has a state $x_v \in Q$
- Dynamic F: If a vertex v with state x_v has a neighbor with state s(x_v) = x_v + 1 mod q then the new state of x_v is s(x_v), if not, the state does not change.

Example

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Some definitions

Cycles:



• Jump J(C) = 3 and Length L(C) = 4

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Period 4

Some definitions 2

Skeleton Es(x): Subgraph of G with the same set of vertices. An edge connect two vertices with state p and p' if and only if p = s(p'), p = p' or s(p) = p'.

- $Es(x) \subseteq Es(F(x))$
- In a finite number of iteration, the skeleton becomes stable.

Example

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Step 0

Step 1

Step 2

Predicting the future

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Proposition 1 (Goles and Matamala). Let x be an assignment and assume that Es(x) is stable. Then for all $L \ge 1$ we have

1. If $a(v, L, x) \leq 0$ then $F_v^l(x) = x_v \ \forall l \in \{1, \dots, L\}$. 2. If a(v, L, x) > 0 then $F_v^l(x) = s^{a(v,L,x)}(x_v) \ \forall l \in \{h(v, L, x), \dots, L\}.$

a(v, L, x) the maximum of the jumps of all the walks of length less or equal than L starting in v on x.

h(v, L, x) be the minimum of the lengths among all walks starting in x which reach the maximum jump a(v, L, x).

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Some definitions 3

 Efficiency of a path P: is the quotient between its jump and its length.

$$e(P, x) = \frac{J(P, x)}{L(P)}$$

• Efficiency of the system or *global efficiency*: is the maximum efficiency over all closed walks in the skeleton.

Previous results [G. & M.]

- Closed walks of global efficiency take control of the dynamic.
- The length of the period divides the least common multiple of the length of the closed walks with global efficiency. [non-polynomial upper bound]
- This upper bound can be reached.

Reaching the upper bound

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To obtain non-polynomial periods in the 2-dimensional lattice \mathbb{Z}^2 .

Complete Skeleton

We can decompose the skeleton in "tiles" (cycles of length 4). Its jumps could be 0, 3 or 4. Property: We can sum the jump of some cycles

$$J(\gamma, x) = J(\gamma_1, x) + J(\gamma_2, x)$$



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Some results

If the skeleton is complete, the efficiency and periods possibles are:

	$q \ge 5$	q = 4	q = 3
e(x)	0	0 or 1	$\geq 3/4$
Periods	1	1 or 4	any even ≥ 18

also

For q=3 and for any $r\in \mathbb{Q}\cup [\frac{3}{4},1]$ exists a Q -assignment x such that e(x)=r

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Our main result

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How to "repeat" the idea of previous construction?

Main problems:

Construction of cycles with a desired length

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- Independent evolution of cycles
- Assure the stability of the skeleton

Steps of the construction

- 1. Good sequences of states
- 2. Construction (iteratively) of a cycle with a desired length
- 3. Completion to a rectangle
- 4. Extension for to embed the construction in the plane

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5. Interaction between different cycles

Good sequences \mathcal{A}

- 1. For every i = 0, ..., kp 1, $a_{i+1 \mod kp} \in \{a_i, s(a_i)\}$ and $a_{i+3 \mod kp} \notin \{a_i, s(a_i)\}$.
- 2. The jump J(a) of the sequence a is kq, where $J(a) := \sum_{i=0}^{kp-2} \mu(a_i, a_{i+1}).$
- 3. For every $j \in \{0, ..., k-1\}$ and for every i = 0, ..., kp - 1 the jump of the subsequence $(x_i, ..., x_{i+jp \mod kp})$ belongs to $\{jq - 1, jq, jq + 1\}$.

Example: (q = 6) 001234450012344500123345

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Main cycles

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• $\forall a \in \mathcal{A}$ the skeleton is the same.

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$$F(x(a)) = x(\sigma(a)).$$

• $Es(G_k, F(x(a))) = Es(G_k, x(a)).$

Completion to a rectangle



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Embeding the construction in the plane



Interaction between cycles

Theorem 1. For every $q \ge 5$ and every even integer psuch that $q and, for every integer <math>m \ge 1$ and integers k_1, k_2, \ldots, k_m there exists a Q-assignment y of K and a vertex v in K such that $T_v(y) = p \cdot \operatorname{lcm}_{i=1,\ldots,m}\{k_i\}.$

<u>Idea:</u> Cycles with sequence a^i embedded equidistant to a particular vertex, where

$$a^{i} = (1, 2, 2, 3, \dots, 0)^{(k_{i}-1)}(1, 1, 2, 3, \dots, 0)$$

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Final Example

- $Q = \{0, 1, 2, 3, 4, 5\}$
- $e(x) = \frac{3}{4}$

• Sequences $(00123445)^{k_i-1}(00123345)$

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- $k_1 = 2$, $k_2 = 3$, $k_3 = 5$
- **Period** = $8 \cdot \text{lcm}(2, 3, 5) = 240$