

Let \mathcal{A} and \mathcal{B} be categories. A functor $F: \mathcal{A} \rightarrow \mathcal{B}$ consists of

(2.6) a function which to each object A of \mathcal{A} assigns an object FA of \mathcal{B} ;

(2.7) a function which to each morphism $f: A_1 \rightarrow A_2$ in \mathcal{A} assigns a morphism $Ff: FA_1 \rightarrow FA_2$ in \mathcal{B} .

The following axioms are postulated.

(2.8) $F(gf) = (Fg)(Ff)$.

(2.9) $F1_A = 1_{FA}$.

Given functors $F: \mathcal{A} \rightarrow \mathcal{B}$ and $G: \mathcal{B} \rightarrow \mathcal{C}$, the composite functor $GF: \mathcal{A} \rightarrow \mathcal{C}$ is defined in the obvious way.

3. EXAMPLES OF CATEGORIES

The category S of sets has sets as objects and functions as morphisms with composition defined as composition of functions. Two objects in S will play special roles: the empty set \emptyset and the set I consisting of the number 1 alone. For any set A there are unique morphisms

$$\emptyset \rightarrow A, \quad A \rightarrow I.$$

We say that \emptyset is an *initial* object and I is a *terminal* object for S .

A set A and an element $a \in A$ determine a unique morphism $I \rightarrow A$ with a as value. We shall denote this morphism by the same letter a . Thus morphisms $I \rightarrow A$ and elements of A will be identified.

For each integer $n = 0, 1, \dots$ we denote by $[n]$ the set $\{1, \dots, n\}$. Thus $[0] = \emptyset$ and $[1] = I$. The sets $[n]$, $n = 0, 1, \dots$ together with all morphisms between them form a subcategory S_0 of S .

4. THEORIES

A *theory* T is a category such that

(4.1) the objects of T are $[n]$ for $n = 0, 1, \dots$;

(4.2) S_0 is a subcategory of T ; i.e., every morphism in S_0 is also a morphism in T , composition of morphisms in S_0 agrees with that in T , and the identity morphisms $1_{[n]}$ in S_0 are also identity morphisms in T ;

(4.3) given morphisms

$$\phi_i: I \rightarrow [p] \text{ in } T, \quad i = 1, \dots, n,$$

there exists a unique morphism

$$\phi: [n] \rightarrow [p]$$

such that ϕ_i is the composition

$$I \xrightarrow{i} [n] \xrightarrow{\phi} [p]$$

for every $i \in [n]$.

We shall write $\langle \phi_1, \dots, \phi_n \rangle$ for the morphism ϕ given in (4.3). Thus for any morphism $\phi: [n] \rightarrow [p]$ in T , we have $\phi = \langle \phi_1, \dots, \phi_n \rangle$.

It should be noted that it follows from the above axioms that in T there is only one morphism $O_n: \emptyset \rightarrow [n]$ for every n , just as in the case of S_0 . However, (contrary to what takes place in S_0) there may be in T morphisms $\phi: I \rightarrow \emptyset$. In fact, these "0-ary operations" play a fundamental role in the sequel.

5. ALGEBRAS

Let T be a theory. A T -algebra A consists of a set A and a rule which to each $\phi: [n] \rightarrow [p]$ in T and each p -tuple (x_1, \dots, x_p) of elements of A assigns an n -tuple

$$(x'_1, \dots, x'_n) = (x_1, \dots, x_p)\phi$$

of elements of A , subject to the following two axioms:

(5.1) if ϕ is in S_0 , then $x'_i = x_{\phi_i}$;

(5.2) if $\psi: [k] \rightarrow [n]$ in T , then

$$(x'_1, \dots, x'_n)\psi = (x_1, \dots, x_p)(\phi\psi).$$

If we write $x = (x_1, \dots, x_p)$, then (5.2) may be rewritten as

$$(x\phi)\psi = x(\phi\psi). \quad (5.2')$$

A morphism $f: A \rightarrow B$ of T -algebras is a mapping from A to B satisfying

$$f[(x_1, \dots, x_p)\phi] = (fx_1, \dots, fx_p)\phi, \quad (5.3)$$

or in abbreviated form

$$f(x\phi) = (fx)\phi, \quad (5.3')$$

where $fx = (fx_1, \dots, fx_p)$.

With composition of morphisms of algebras defined in the ordinary fashion, there results the category T^b of T -algebras.

We note that if $\phi: I \rightarrow [p]$ in T , then

$$(x_1, \dots, x_p)\phi \in A$$

so that ϕ yields a mapping $A^p \rightarrow A$ where A^p is the p -fold Cartesian product $A \times \dots \times A$.

If in the above $p = 0$, i.e., $\phi: I \rightarrow \emptyset$, then $()\phi \in A$ is an element of A determined by ϕ , independent of any "inputs" a_1, \dots, a_p . We denote this element by ϕ_A .