Smooth manifold reconstruction from point clouds

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From unorganized 3D point clouds to triangulated surfaces : how to connect the dots ?



Overview

- Preliminaries : normal estimation and poles
- Combinatorial methods
- Implicit methods
- Continuation methods
- Multi-scale approaches

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Estimating normal directions



Voronoi regions are elongated in the normal direction

 $V(p_i)$ contains the center m_{p_i} of the medial ball tangent to *S* at p_i

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Pole of p_i : vertex v of $V(p_i)$ furthest from p_i

 $\|\boldsymbol{v}-\boldsymbol{p}_i\| \geq \|\boldsymbol{p}_i - \boldsymbol{m}_{\boldsymbol{p}_i}\| \geq \mathrm{lfs}(\boldsymbol{x})$

Error bound on the normal direction

 $\mathcal{P} = \varepsilon$ -sample v pole of $x : ||v - x|| \ge lfs(x)$ $\|i - x\| \le \varepsilon \operatorname{lfs}(i) \le \frac{\varepsilon}{1-\varepsilon} \operatorname{lfs}(x)$ (lfs 1-Lipschitz) $\angle(\vec{n}(\mathbf{x}), \vec{\mathbf{vx}}) = \alpha + \beta$ $\leq 2 \arcsin \frac{\varepsilon}{1-\varepsilon}$ $\approx 2\varepsilon$



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Orientation of the normals

orientation may be given by the scanning device

otherwise :



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- orient the normals at the vertices of $conv(\mathcal{P})$ (easy)
- then propagate the labels coherently
 - p_i close to $p_j \Rightarrow n_{p_i} \cdot n_{p_i} > 0$
 - heuristics : walk along the MST of \mathcal{P} [Hoppe 92]

Warning !

The problem is global : local methods are non robust in the presence of

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Crust and power crust

[Amenta et al.]

Noise, sharp features and undersampling

- Spectral surface reconstruction
- Voronoi PCA

[Kolluri et al. 2006] [Alliez et al. 2007] [Mérigot 2009]

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Power crust algorithm

[Amenta et al.]



- 1. For each $p_i \in \mathcal{P}$, compute
 - its pole m_{p_i} and $w_i = ||m_{p_i} p_i||$
 - ► the secondary pole : $m'_{p_i} = \arg \min_{v \in V(p_i)} ((m_{p_i} - p_i) \cdot (m'_{p_i} - p_i))$ and $w'_i = ||m'_{p_i} - p_i||$
- 2. Compute the weighted Voronoi (power) diagram WVor(WP) where $WP = \bigcup_{i=1..n} \{(m_{p_i}, w_i)\} \bigcup \{(m'_{p_i}, w'_i)\}$
- 3. Label the elements of WP as inside or outside
- Output the facets of WVor(WP) incident to cells with opposite labels

Power crust algorithm

[Amenta et al.]



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Global labelling of poles using a spectral approach

Kolluri, Shewchuk, O'Brien Spectral surface reconstruction from noisy point clouds Siggraph 2004

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[Kolluri, Shewchuk, O'Brien]

Surface reconstruction : the implicit approach

- Compute a function over ℝ³ whose zero-set Z either interpolates or approximates E
- 2. mesh Z



Various implicit functions

- interpolation of the sign distance functions to the tangent planes at the sample points [Hoppe 92], [B. & Cazals 00]
- moving least squares (MLS)
- radial basis functions
- Poisson based approach
- spectral method

[Levin 03]

- [Carr et al. 01]
- [Khazdan et al. 06]
 - [Alliez et al. 07]



$$h_i(x) = (x - p_i) \cdot n_i$$

assuming we know or can estimate the normals n_i (direction and orientation) to S at p_i

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Functional interpolant $h(x) = \sum_{i} \lambda_{i}(x)h_{i}(x)$ where $\forall x, \sum_{i=1}^{n} \lambda_{i}(x) = 1$ (partition of unity)



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Implicit surface $h^{-1}(0)$



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Implicit surface $h^{-1}(0)$

Properties

- $\lambda_i(p_j) = \delta_{ij} \Rightarrow h(p_i) = h_i(p_i) = 0$
- ► if the λ_i are C^k continuous, h⁻¹(c) is a C^k surface for almost all c (by Sard's th.)

Choice of $\lambda_i(x)$

Some examples

- Nearest neighbor (NN) : λ_i(x) = 1, λ_{j≠i}(x) = 0 if x ∈ V(p_i) Not continuous, support = V(p_i)
- ► Radial basis function : λ_i(x) = exp^{-β||x-p_i||²}, β > 0 C[∞], non compact support
- Natural neighbor coordinates
 C¹-continuous, compact support

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Natural Neighbor Coordinates



[Sibson 80]

Natural neighbors

For $x \in \operatorname{conv}(E)$ Nat $(x) = \{p_i : w_i(x) \neq 0\}$

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Natural Neighbor Coordinates



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Laplace's coordinates

$$\lambda_i(\mathbf{x}) = \frac{v_i(\mathbf{x})}{\sum_i v_i(\mathbf{x})}$$
 with $v_i(\mathbf{x}) = \frac{\operatorname{vol}(v^+(\mathbf{x}, p_i))}{\|\mathbf{x} - p_i\|}$, $\lambda_i(p_i) = 1$, $\lambda_j(p_i) = 0$

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Sibson's coordinates

$$\sigma_i(x) = \frac{w_i(x)}{\sum_i w_i(x)}$$
 with $w_i(x) = \operatorname{vol}(V^*(p_i) \cap V^*(x))$

Equivalently $Nat(x) = \{ vertices of star(x) in Del^+(E) \}$ = { vertices of the Delaunay simplexes whose circumscribing ball $\ni x \}$



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 \mathbb{R}^d /Nat is the arrangement of the Delaunay spheres

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 \mathbb{R}^d /Nat is the arrangement of the Delaunay spheres

Support of λ_i or σ_i = union of the Delaunay balls incident to p_i

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Properties of the Laplace coordinates $\lambda_i(x)$ $\lambda_i(x) = \frac{v_i(x)}{\sum_i v_i(x)}$ with $v_i(x) = \frac{\text{vol}(V^+(x,p_i))}{\|x-p_i\|}$

- 1. Trivially $\sum_i \lambda_i(x) = 1$
- 2. The λ_i are continuous

3. when
$$x \to p_i$$
, $v_i(x) \to \infty$
 $v_j(x) \to \operatorname{cst} < \infty$ (j
 \Rightarrow $\lambda_i(x)) \to 1$
 $\lambda_j(x)) \to 0$



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4.
$$x = \sum_i \lambda_i(x) p_i$$
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$$x = \sum_i \lambda_i(x) p_i$$

Minkowski's th.



$$d \operatorname{vol}(P) = \sum_{j \in J} \operatorname{vol}(f_j) (x - p_j) \cdot n_j$$

 $\nabla \operatorname{vol}(P) = 0 \Rightarrow \sum_{j \in J} \operatorname{vol}(f_j) n_j = 0$

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Applied to $V^+(x)$

Hence
$$\sum_{i} v_i \frac{x - p_i}{\|x - p_i\|} = 0 \iff \left(\sum_{i} \frac{v_i}{\|x - p_i\|}\right) x = \sum_{i} \frac{v_i}{\|x - p_i\|} p_i$$

Properties of the Sibson's coordinates $\sigma_i(x)$

1. Trivially
$$\sum_{i} \sigma_i(x) = 1$$

2. The σ_i are continuous

3. when
$$x \rightarrow p_i$$
, $w_i(x) \rightarrow C \neq 0$
 $w_j(x) \rightarrow 0$ $(j \neq i)$
 \Rightarrow $\sigma_i(x)) \rightarrow 1$
 $\sigma_j(x)) \rightarrow 0$



4.
$$x = \sum_i \sigma_i(x) p_i$$

same proof as for the Laplace coord. but in the space of spheres

$\mathbf{x} = \sum_i \sigma_i(\mathbf{x}) \mathbf{p}_i$

Minkowski's th. applied to $\mathcal{V}_p = h_x^- \cap \{h_i^+\}$



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$$\mathbf{x} = \sum_i \sigma_i(\mathbf{x}) \mathbf{p}_i$$

Minkowski's th. applied to $\mathcal{V}_p = h_x^- \cap \{h_i^+\}$



$$n_{i} = (2 p_{i}, -1) \perp f_{i}$$

$$n_{p} = (-2 x, 1) \perp f_{x}$$

$$\sum_{i} \operatorname{vol}(f_{i}) \frac{n_{i}}{\|n_{i}\|} + \operatorname{vol}(f_{x}) \frac{n_{x}}{\|n_{x}\|} = 0$$

$$w_{i}(x) = \operatorname{vol}(V^{+}(p) \cap V(p_{i})) =$$

$$\operatorname{vol}(f_{i}) \frac{n_{i}}{\|n_{i}\|} \cdot (-i_{d+1}) = \frac{\operatorname{vol}(f_{i})}{\|n_{i}\|}$$

$$w(x) = \operatorname{vol}(V^{+}(x)) = \frac{\operatorname{vol}(f_{x})}{\|n_{x}\|}$$

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$$\sum_{i} w_i(x)n_i + w(x)n_x = 0 \implies \sum_{i} w_i(x)p_i - w(x)x = 0$$

Strength of the implicit approach : Refining the mesh allows to remove singularities





(Using Sibson's interpolation)

Certified algorithms and codes

- ► Cocone algorithm [Dey & al.]
 - Geomagic Studio
 - Natural neighbor interpolation of distance functions (Catia)
 - Numerical approaches [Kolluri & al., Khazdan & al., Alliez & al.]
 - CGAL-3.5 :

Geometry Processing & Surface Reconstruction [Alliez & al.]

Smooth submanifold reconstruction

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- Motivation and difficulties
- Continuation approach and intrinsic complexity
- Dimension estimation and multi-scale reconstruction

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Motivations

- Theoretical foundations and guarantees (Sampling theory for geometric objects)
- Complexity issues and practical efficiency
- Higher dimensions
 - Reconstruction in 3D+t space
 - 6D phase space
 - Configuration spaces of robots, molecules...
 - Data analysis

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The curses of Delaunay triangulations in higher dimensions

- Subdividing the ambient space is very costly
 - ? Can we store $Del(\mathcal{P})$ in a more compact way?
 - ? Can we construct more local data structures?
- The restricted Delaunay triangulation is no longer a good approximation of the manifold even under strong sampling conditions
 - ? What else ?

Is Delaunay triangulation of practical use in high dimensional spaces?

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Combinatorial complexity of the Delaunay triangulation

• Worst-case: $O(n^{\left\lceil \frac{d}{2} \right\rceil})$

• Uniform distribution in \mathbb{R}^d : $c^d O(n)$

[Dwyer 1991]

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Deciding whether $x \in B(p_0, ..., p_d)$ reduces to evaluating the sign of

$$\sigma(p_0, ..., p_d, x) = \begin{vmatrix} 1 & ... & 1 & 1 \\ p_0 & ... & p_d & x \\ p_0^2 & ... & p_d^2 & x^2 \end{vmatrix}$$

 σ is a polynomial of degree d + 2 in the input variables (its exact evaluation would require (d + 2)-fold precision)

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Implementation of the incremental algorithm [Hornus 2009]

- explicit storage of the d-simplices
- and of their adjacency graph



- Exact computing
- To come in CGAL 3.6

Compact representation

Delaunay graph

Store only the Delaunay graph (edges + vertices)
 + 1 *d*-simplex per vertex σ(p)

▶ insert(X):

- 1. locate the vertex of $Del(\mathcal{P})$ nearest to x
- remove all *d*-simplices whose circ. ball ∋ x by walking from neighbor to neighbor in Del(P)
- 3. update the Delaunay graph by joining *x* to the vertices of the removed simplices



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Performances of the Delaunay graph construction

1	Dimension	2	3	4	5	6
2	Number of input points	1024K	1024K	1024K	256K	32K
3	Size of the simplex-cache	1K	1K	10K	300K	1000K
4	Size of the conflict zone	4.1	21	134	940	6145
5	Calls to neighbor(,)	12.2	84.6	671.2	5631	43021
6	Number of candidates	2	2.6	4	6.7	11.6
7	Fast cache hit (non- <i>null</i> pointer)	56.6%	57.5%	54.6%	55.5%	54.3%
8	Cache hit	37%	39.6%	40.1%	42.3 %	43.1 %
9	Cache miss	6.4 %	2.9%	5.3%	2.2 %	2.6%
10	Time ratio (Del_graph/New_DT)	6.1	5.7	6.0	6.5	8.1
11	Space ratio (Del_graph/New_DT)	2.7	1.7	0.6	0.2	0.1
12	Number of simplices per vertex	6	27(×4.5)	$157(\times 5.8)$	1043(×6.7)	7111(×6.8)
13	Number of edges per vertex	6	$15.5(\times 2.6)$	$36.5\scriptscriptstyle (\times 2.4)$	73(×2)	$164.6(\times 2.25)$

Current best: 100K 6D vertices: 105 millions simplices. 15 hours for the graph (approx 2:30 hours for full-D).

Space can be further improved using a compact representation of graphs [Blandford et al 2003] Is the restricted Delaunay triangulation a good approximation

in high dimensional spaces?

Winter School on Algorithmic Geometry Smooth manifold reconstruction

Tangent space approximation

 \mathbb{M} is a smooth *k*-dimensional manifold embedded in \mathbb{R}^d

Bad news

The Delaunay triangulation restricted to \mathbb{M} may be a bad approximation of the manifold even if the sample is dense



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Tangent space approximation

 \mathbb{M} is a smooth *k*-dimensional manifold embedded in \mathbb{R}^d

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The Delaunay triangulation restricted to \mathbb{M} may be a bad approximation of the manifold even if the sample is dense



Good news

[Cheng et al. 2005]

If τ nor its faces are slivers, there exist a constant a_k (depending on σ_0) s. t. sin $\angle(\operatorname{aff}(\tau), T_p) \le a_k \varepsilon$

Slivers in higher dimensions

Sliver

Given constants ρ_0 and σ_0 , a *j*-simplex τ is called a sliver if

1.
$$j > 2$$

2. $\rho_{\tau} = \frac{R_{\tau}}{L_{\tau}} \le \rho_{0}$
3. $\sigma_{\tau} = \operatorname{vol}(\tau)/\mathcal{L}_{\tau}^{j} < \sigma_{0}$
4. $\forall \sigma \subset \tau, \, \rho_{\sigma} \le \rho_{0} \text{ and } \sigma_{\sigma} \ge \sigma_{0}$

$$m{R}_{ au} = {
m radius} \; {
m of} \; {
m the} \; {
m circ.} \; {
m ball} \ m{L}_{ au} = {
m length} \; {
m of} \; {
m the} \; {
m shortest} \; {
m edge} \; {
m of} \; au$$



Slivers are fragile



A small perturbation of one of its vertices may lead to

a negative volume

 $vol(\tau)$ is small

a big circumscribing sphere

In both cases, the sliver is removed from DT

Two certified techniques to remove slivers

Weight the vertices, replace the Delaunay triangulation by the regular triangulation of the WP [Cheng et al. 2000]

- the points are not moved
- must be used as a postprocessing step on the mesh
- Perturbe the vertices

[Li 2001]

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can be used during Delaunay refinement

Weighted Delaunay triangulation

Weight assignment

$$p_i \rightarrow \omega(p_i), \qquad \omega = (\omega(p_1), ..., \omega(p_n))$$

Weighted Voronoi diagram

$$V^{\omega}(p_i) = \{x : \|x - p_i\|^2 - \omega^2(p_i) \le \|x - p_j\|^2 - \omega^2(p_j)\}$$

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Weighted Delaunay triangulation

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Weighting a vertex





A continuation approach to manifold reconstruction

- Can we compute $\operatorname{Del}_{S}^{\omega}(\mathcal{P})$ without computing $\operatorname{Del}(\mathcal{P})$?
- Can we avoid subdividing the embedding space and obtained an intrinsic complexity ?

Basic assumption in manifold learning

Data live in a low-dimensional manifold embedded in a high-dimensional space

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Example 1 : human face images



head = sphere camera : 3 dof light source : 2 dof

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An image with *N* pixels \rightarrow a point in \mathbb{R}^N

It is impossible to triangulate points in such a huge space !

Example 1 : human face images



head = sphere camera : 3 dof light source : 2 dof

An image with N pixels \rightarrow a point in \mathbb{R}^N

It is impossible to triangulate points in such a huge space !

Example 2 : points with unit normals

•
$$(p_i, \mathbf{n}_i) \in \mathcal{N} = \mathbb{R}^3 \times \mathbb{S}^2$$

 The surface to be reconstructed is a 2-manifold embedded in N

The tangential Delaunay complex

[Freedman 2002], [B.& Flottoto 2004], [B.& Ghosh 2009]

[Cheng, Dey, Ramos 2005]



- Construct the star of p ∈ P in the Delaunay triangulation Del_{Tp}(P) of P restricted to T_p
- 2. $\operatorname{Del}_{TM}(\mathcal{P})$: the set of stars of p, $p \in \mathcal{P}$

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- + $\operatorname{Del}_{\mathcal{T}\mathbb{M}}(\mathcal{P}) \subset \operatorname{Del}(\mathcal{P})$
- + star(*p*), Del_{T_p}(*P*) and therefore Del_{TM}(*P*) can be computed without computing Del(*P*)
- $\mathrm{Del}_{\mathcal{TM}}(\mathcal{P})$ is not necessarily a triangulated manifold

Construction of $\text{Del}_{T_p}(\mathcal{P})$

Given a *k*-flat *H*, $Vor(\mathcal{P}) \cap H$ is a weighted Voronoi diagram



$$\begin{aligned} \|x - p_i\|^2 &\leq \|x - p_j\|^2 \\ \Leftrightarrow \quad \|x - p_i'\|^2 - \|p_i - p_i'\|^2 &\leq \|x - p_i'\|^2 - \|p_j - p_j'\|^2 \end{aligned}$$

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Construction of $\text{Del}_{T_p}(\mathcal{P})$

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Corollary: construction of Del_{T_p}

- 1. project \mathcal{P} onto T_{ρ} which requires O(dn) time
- 2. construct star(p'_i) in $\text{Del}^{\omega}(p'_i) \subset T_{p_i}$ where $\omega(p_i) = \|p_i p'_i\|$ 3. star(p_i) $\stackrel{1-1}{\leftrightarrow}$ star(p'_i)

Inconsistent configurations $\phi = [p_1, ..., p_{k+2}]$

Definition $\exists p_i, p_j, p_l \in \phi \text{ s.t.}$

- 1. $\tau = \phi \setminus \{p_l\} \in \operatorname{star}(p_i)$ $\notin \operatorname{star}(p_j)$
- 2. τ nor its faces are slivers
- Vor(*p*_l) is the first Voronoi cell intersected by *c*_i*c*_j



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Observations

• c_i and c_j are close if p_i is close to p_j and $N(\tau) \approx N(p_i)$ $\Rightarrow \phi$ is a (k + 1)-sliver

• $\phi \in \text{Del}(\mathcal{P})$

Inconsistency removal by weighting \mathcal{P}

 $\omega = \text{weight assignment}$

 $I\!F^\omega$: set of faces of the inconsistent configurations

1. For
$$j = 1..k$$
, for $i = 1..n$

weight p_i so as to remove all slivers incident to p_i that are in $\text{Del}^{\omega}_{TM}(\mathcal{P})$ and IF^{ω}

2. for *i* = 1..*n*

weight p_i so as to remove all inconsistent configurations

Basic operations

- ► Compute the tangent space at a point of M
- Project a point on a k-flat
- Maintain the star of a point when varying its weight
- No d-dimensional data structure

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Basic operations

- Compute the tangent space at a point of \mathbb{M}
- Project a point on a k-flat
- Maintain the star of a point when varying its weight
- ► No *d*-dimensional data structure

Properties of the output

If \mathcal{P} is a sparse (not necessarily uniform) ε -sample, upon termination, the stars are coherent, the simplices are small and they locally approximate the tangent space of \mathbb{M}

- $Del_{TM}(\mathcal{P})$ is a PL simplicial *k*-manifold
- \blacktriangleright isotopic and close to $\mathbb M$

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Dimension of S?



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Witness complex and multiscale reconstruction

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Relaxing the definition of restricted Delaunay triangulation [Carlsson & de Silva 2004]



- witnesses : $W \subset \mathbb{M}$ (not necessarily finite)
- landmarks : a finite set of points L ⊂ W
- τ is a weak Delaunay simplex iff
 - $\tau \subset L$
 - $\forall \sigma \subseteq \tau, \exists w \in W$ closer to σ than to $L \setminus \sigma$,

Weak Delaunay (witness) complex Wit(L, W)

the collection of all weak Delaunay simplices σ , i.e. σ and all its faces have a witness in *W* with respect to *L*



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Weak Delaunay (witness) complex Wit(L, W)

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Easy to compute (only distance comparisons)

```
Clearly, \text{Del}(L) \subset \text{Wit}(L, \mathbb{R}^d)
```

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WW-Theorem : $\forall \Omega \subset \mathbb{R}^d$, Wit $(L, \Omega) \subset$ Wit $(L, \mathbb{R}^d) \subset$ Del(L)

Corollaries

- Wit $(L, \mathbb{R}^d) = \text{Del}(L)$
- Wit(L, Ω) is embedded in \mathbb{R}^d

Proof of the Weak Witness theorem

 $\tau = [p_0, ..., p_k]$ is a *k*-simplex of Wit(*L*) witnessed by a ball W_{τ} , i.e. $W_{\tau} \cap L = \tau$

We prove that $\tau \in \text{Del}(L)$ by a double induction on

$$k \\ |S_{\tau} \cap \tau| \qquad (S_{\tau} = \partial W_{\tau})$$

Clearly true for k = 0 and $|S_{\tau} \cap \tau| = k + 1$

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Proof of the weak witness theorem





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Case of sampled domains : bad news

 \mathcal{P} a finite set of points $\subset \Omega$, $L \subset \mathcal{P}$

Wit(L, \mathcal{P}) \neq Del(L, Ω), even if \mathcal{P} is a dense sample of Ω



 $[ab] \in \operatorname{Wit}(L, \mathcal{P}) \iff \exists p \in \mathcal{P}, \operatorname{Vor}_2(a, b) \cap \mathcal{P} \neq \emptyset$

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Case of sampled surfaces : good news

 $\mathcal P$ a uniform $\varepsilon\text{-sample}$ of a surface $\mathcal S\subset\mathbb R^3$

 \mathcal{P}^{ξ} the set of balls $B(p, \xi \operatorname{lfs}(\mathcal{S})), p \in \mathcal{P}$

if

 $L \subset \mathcal{P}$ a uniform λ -sample of \mathcal{P} (landmarks)

Although Wit(L, \mathcal{P}^{ξ}) may not be a triangulated surface, all its facets are close (both in position and orientation) to S, which makes surface extraction easy

Case of sampled manifolds : good news

[B.,Guibas, Oudot 2008]

 \mathcal{P} a uniform ε -sample of a surface $\mathcal{S} \subset \mathbb{R}^3$

```
\mathcal{P}^{\xi} the set of balls B(p, \xi \operatorname{lfs}(\mathcal{S})), p \in \mathcal{P}
```

 $L \subset \mathcal{P}$ a uniform λ -sample of \mathcal{P} (landmarks)

 $\exists \xi_1, \xi_2$ and a weight assignment ω s.t.

By weighting the points of *L*, we can remove slivers from $Wit(L, \mathcal{P}^{\xi_2})$ and obtain a triangulated manifold $\approx S$

Theorem

[B., Guibas, Oudot 2007]

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If $\xi_1 \geq \varepsilon$ and $30\varepsilon \leq \lambda \ll lfs(S)$, one can compute a weight assignment ω s.t. Wit^{ω}(L, \mathcal{P}^{ξ_1}) = Del^{ω}(L, S) $\approx S$

Theorem

[B., Guibas, Oudot 2007]

If $\xi_1 \geq \varepsilon$ and $30\varepsilon \leq \lambda \ll lfs(S)$, one can compute a weight assignment ω s.t. Wit^{ω}(L, \mathcal{P}^{ξ_1}) = Del^{ω}(L, S) $\approx S$

Multiscale reconstruction

- In practice, ε and lfs(S) are unknown
- There may exist S₁...S_i s.t. W is an ε_i-sample of S_i



Theorem

[B., Guibas, Oudot 2007]

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Multiscale reconstruction

- In practice, ε and lfs(S) are unknown
- There may exist S₁...S_i s.t. W is an ε_i-sample of S_i



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- Generate a monotonic sequence of samples L ⊆ W
- As long as $30\varepsilon_i \leq \lambda \leq lfs(S_{i+1})$, $Wit^{\omega}(L, \mathcal{P}^{\xi}) \approx S_i$
- ► This can be detected by looking for plateaus in the diagram of Betti numbers of Wit^ω(L, W^ξ)

Manifold reconstruction algorithm

Greedy max min algorithm maintain Wit^{ω}(L, \mathcal{P}^{ξ_1}) and Wit^{ω}(L, \mathcal{P}^{ξ_2}) $\xi_1 = \lambda/30, \xi_2 = 3\lambda$ INIT L := a point of \mathcal{P} REPEAT $L \leftarrow p$ = the point of $\mathcal{P} \setminus L$ furthest from Lcompute $\omega(p)$ so as to remove slivers in Wit^{ω}(L, \mathcal{P}^{ξ_2}) update Wit(L, \mathcal{P}^{ξ_1}) and Wit(L, \mathcal{P}^{ξ_2}) UNTIL $L = \mathcal{P}$



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Manifold reconstruction algorithm





Update Wit(L, p^{ξ})

maintain the *k*-order Vor(*L*), $k \le d + 1$

(The curse of dim. is back)

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 $[p_1...p_k] \in \operatorname{Wit}(L, \mathcal{P}^{\xi}) \Leftrightarrow \exists w \in \mathcal{P}, \operatorname{Vor}_k(p_1...p_k) \cap B(p, \xi) \neq \emptyset$

Multi-scale reconstruction

- maintain the Betti numbers of \hat{S} on the fly
- detect the plateaus
- return the corresponding Ŝ that are valid approximations of S (given the sample P) at different scales

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Multi-scale reconstruction

- maintain the Betti numbers of \hat{S} on the fly
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Remarks

- *P* is not required to be sparse
- lfs(S) needs not to be known

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